The alternatives of bare and modified numerals

3 (BNS) more/less than 3 (CMNs) (ScalAlts), (SubDomAlts) at most/least 3 (SMNs)

(ScalAlts) (ScalAlts), SubDomAlts

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Preview

- * 3, more/less than 3, and at least/most 3 differ w.r.t. (at least)
 - entailments,
 - scalar implicatures,
 - ignorance, and
 - acceptability in downward-entailing environments.
- ★ Many theories have been proposed to capture these differences.
- ★ Lately a move towards alternative-based theories.
- ★ Promising results, but also empirical and conceptual issues.
- ★ I will propose a theory that overcomes these issues.

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Entailments

- * 3 / more than 3 / at least 3 carry lower-bounding entailments.
- (1) a. Alice has 3 diamonds.
 b. → not 2 or less
 c. Alice has 3 diamonds, # if not less.
- ★ less than 3 / at most 3 carry upper-bounding entailments.
- (2) a. Alice has less than 3 diamonds.
 - b. --> not 3 or more
 - c. Alice has less than 3 diamonds, #if not more.

★ Existing proposals: Multiple possible solutions, typically not compositional down to the smallest pieces.

* We want one that gets these entailments with ease and also minimally uncovers the uniform contribution of the numeral, *much/little*, or [-er]/[at -est] in producing these entailments.

Scalar implicatures I

- ★ BNs also carry upper-bounding scalar implicatures. [Horn, 1972]
- (3) a. Alice has 3 diamonds.
 b. → not 4 yields 'exactly 3' meaning ✓
 c. Alice has 3 diamonds, if not 4.

- ★ CMNs and SMNs don't seem to. [Krifka, 1999]
- (4) a. Alice has more than 3 diamonds.
 b. ৵ not more than 4 yields 'exactly 4' meaning X

★ Existing proposals: No scalar implicatures for CMNs and SMNs.

Scalar implicatures II

- ★ But in certain contexts all give rise to scalar implicatures!
- (5) a. If you have at least 3 diamonds, you win.b. → not if at least 2

- ★ And in some none do:
- (6) a. Alice doesn't have 3 diamonds.
 b. ≁ not not 2 yields 'exactly 2' meaning X

- * We want scalar implicatures for all!
- \star We need a separate mechanism to rule out certain implicatures.

Scalar implicatures III

 ★ With coarser granularity, CMNs and SMNs can give rise scalar implicatures too. [Spector, 2014, Cummins et al., 2012, Enguehard, 2018]

- (7) Grades are given based on the number of problems solved. People who solve more than 5 problems but fewer than 9 problems get a B, and people who solve 9 problems or more get an A.
 a. John solved more than 5 problems.
 b. w not more than 9 (he gets a B) example from [Spector, 2014]
- ★ That is true of BNs in the problem cases also.
- (8) a. Alice doesn't have 3 diamonds.
 b. ≁ not not 1 (she does have some)

Ignorance I

- ★ SMNs give rise to strong speaker ignorance inferences.
- (9) I have 3 / more than 2 / ??at least 3 children.

★ Existing proposals: e.g., [Büring, 2008, Kennedy, 2015, Spector, 2015]

- SMNs are underlyingly disjunctive (at least 3 = exactly 3 or more than 3) and have domain alternatives (the individual disjuncts).

- Ignorance inferences are implicatures from these alternatives.
- Nothing of this sort is assumed / derived for CMNs.

Ignorance II

- CMNs give rise to ignorance inferences too! [Cremers et al., 2017]
 (10) [A:] How many diamonds does Alice have? [B:] More than 3.
- ★ Unlike BNs and like SMNs, CMNs are compatible with ignorance:
- (11) I don't know how many diamonds Alice has, but she has # 3 / more than 3 / at least 3.
- ★ Unlike CMNs, SMNs are incompatible with exact knowledge.

[Nouwen, 2015]

(12) There were exactly 62 mistakes in the manuscript, so that's more than 50 / # at least 50.

- * We want ignorance implicatures for CMNs too!
- \star We want ignorance to be weaker for CMNs than for SMNs.

Acceptability in DE environments I

★ SMNs are bad under negation.

[Nilsen, 2007, Geurts and Nouwen, 2007, Cohen and Krifka, 2014, Spector, 2015]

(13) Alice doesn't have *at least three / *at most three diamonds. \rightarrow Alice has 2 or less / 4 or more diamonds.

 ★ Existing proposals: The domain alternatives of SMNs are obligatory and must lead to a stronger meaning, but that cannot happen in a DE environment like negation. [Spector, 2015]

Acceptability in DE environments II

★ SMNs are okay in the antecedent of a conditional or the restriction of a universal!

[Geurts and Nouwen, 2007, Cohen and Krifka, 2014, Spector, 2015]

(14) If Alice has at least 3 diamonds, she wins.

(15) Everyone who has at least 3 diamonds wins.

★ We want a solution that can distinguish between various types of DE environments!

Summary and preview of proposal

- BNs, CMNs, and SMNs are non-uniform w.r.t. Entailments Scalar implicatures Ignorance Acceptability in DE environments
- ★ The existing alternative-based proposals are promising, but still:
 - they take into evidence an incomplete dataset;
 - they make non-uniform stipulations about the alternatives;
 - they fail to capture all the patterns we saw.
- \star In this talk:
 - I take into evidence a revised and extended dataset;

- I derive the alternatives of BNs, CMNs, and SMNs in a uniform way from their truth conditions;

- I show how, with certain general assumptions about implicature calculation, we get all the patterns we saw.

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Proposal: Truth conditions and presupposition

the numeral	[Link, 1983, Buccola and Spector, 2016]
$\llbracket n \rrbracket = n$	$\llbracket is_{Card} \rrbracket (n) = \lambda x_e . x = n$
much/little	[Seuren, 1984, Kennedy, 1997]
$\llbracket \text{much} \rrbracket (n) = \lambda d . d \leq n$	$\llbracket \text{little} \rrbracket (n) = \lambda d . d \ge n$
truth conditions 🕕 [Krifka, 1999, Von Stechow, 2005, Heim, 2007, Hackl, 2009]	
(∃ (n P))(Q) [comp](much/little)(n)(P)(Q) [at-sup](much/little)(n)(P)(Q)	$= 1 \text{ iff } \exists x[x = n \land P(x) \land Q(x)]$ = 1 iff $ P \cap Q \in \overline{[[much/little]](n)}$ = 1 iff $ P \cap Q \in [[much/little]](n)$
the presupposition of at-sup	[Hackl, 2009, Gajewski, 2010]

 $| [[much/little]](n)| \ge 2$

✓ Entailments

(16)
$$3 P Q$$
:
 $\exists x[|x| = 3 \land P(x) \land Q(x)]$ (1.b.)

(17) more than 3 P Q:

$$|P \cap Q| \in \overline{[[much]]}(3) \Leftrightarrow |P \cap Q| \in \{4, 5, ...\}$$
 (l.b.)

(18) less than
$$3 P Q$$
:
 $|P \cap Q| \in \overline{[[little]]}(3) \Leftrightarrow |P \cap Q| \in \{\dots, 0, 1, 2\}$ (u.b.)

(19) at most 3 P Q:

$$|P \cap Q| \in \llbracket \text{much} \rrbracket (3) \Leftrightarrow |P \cap Q| \in \{\dots, 0, 1, 2, 3\}$$
(u.b.)

(20) at least 3 P Q:

$$|P \cap Q| \in [[little]] (3) \Leftrightarrow |P \cap Q| \in \{3, 4, ...\}$$
 (l.b.)

Proposal: Alternatives

Scalar alternatives: Replace the *n*-domain with an *m*-domain.

BNs: $\{\exists x[|x| = m \land P(x) \land Q(x)] : m \in S\}$

CMs: { $|P \cap Q| \in [[much/little]](m) : m \in S$ }

SMs: { $|P \cap Q| \in [[much/little]](m): m \in S$ }

Subdomain alternatives: Replace the *n*-domain with its subsets.

BNs: NA (the numeral argument is just a degree) CMs: $\{|P \cap Q| \in A : A \subseteq \overline{[much/little]](n)}\}$ SMs: $\{|P \cap Q| \in A : A \subseteq [[much/little]](n)\}$ active by presup!

obligatory exhaustification relative to SubDomAlts

Examples

(21) BNs: 3 P Q

a. Truth conditions: $\exists x[|x| = 3 \land P(x) \land Q(x)]$ b. ScalAlts: {..., $\exists x[|x| = 2...], \exists x[|x| = 4..., ...}$ c. SubDomAlts: NA

(22) CMNs: e.g., more than 3 P Q

a. Truth conditions: $|P \cap Q| \in \llbracket \text{much} \rrbracket$ (3) b. ScalAlts: {..., $|P \cap Q| \in \llbracket \text{much} \rrbracket$ (2), $|P \cap Q| \in \llbracket \text{much} \rrbracket$ (4), ...}

c. SubDomAlts: $\{|P \cap Q| \in A : A \subseteq [[much]] (3)\}$

(23) SMNs: e.g., at least 3 P Q

a. Truth conditions: $|P \cap Q| \in [[little]]$ (3)b. ScalAlts: {..., $|P \cap Q| \in [[little]]$ (2), $|P \cap Q| \in [[little]]$ (4), ...}c. SubDomAlts: { $|P \cap Q| \in A : A \subseteq [[little]]$ (3)}

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Proposal: Implicature calculation system [Chierchia, 2013]

to exhaustify the scalar alternatives of BNs, CMNs, and SMNs

$$(24) \llbracket O_{ALT}(p) \rrbracket = p \land \forall q \in ALT \llbracket q \to p \subseteq q \rrbracket$$

^S to exhaustify the subdomain alternatives of CMNs and SMNs

- ★ A version of *O* that
 - takes into account presuppositions:

(25)
$$\left[\!\left[O_{ALT}^{S}(p)\right]\!\right] = \pi(p) \land \forall q \in ALT \left[\pi(q) \to \pi(p) \subseteq \pi(q)\right],$$

- requires a properly stronger result:

(26) $[\![O_{ALT}^{PS}(p)]\!]$ is defined iff $O_{ALT}^{S}(p) \subset p$. Whenever defined, $[\![O_{ALT}^{PS}(p)]\!] = [\![O_{ALT}^{S}(p)]\!]$.

last resort, silent, matrix-level, universal doxastic modal

Implicatures from ScalAlts: Scalar implicatures

(27) Alice has 3 diamonds.

a.
$$O_{ScalAlts} (\exists x [|x| = 3 \land P(x) \land Q(x)] \land)$$

= $\exists x [|x| = 3 \land P(x) \land Q(x)] \land$
 $\neg \exists x [|x| = 4 \land P(x) \land Q(x)] \land ...$

(28) Alice has at least 3 diamonds.

a.
$$O_{ScalAlts} (|P \cap Q|) \in [[little]] (3))$$

= $|P \cap Q| \in [[little]] (3) \land$
 $\neg |P \cap Q| \in [[little]] (5) \land ...$

not 4 🗸



(29) If Alice has more than 3 diamonds, she wins.

a. $O_{ScalAlts} ([|P \cap Q|] \in \overline{[much]](3)} \rightarrow win]$ = $[|P \cap Q| \in \overline{[much]](3)} \rightarrow win] \land$ $\neg [|P \cap Q| \in \overline{[much]](2)} \rightarrow win] \land \dots$ not if more than $2\checkmark$

- ★ And so on. We can derive all the attested scalar implicatures.
- ★ Scalar implicatures may be restricted by granularity.
- * In unembeded contexts this effect is compounded by ignorance. 19/38

Implicatures from SubDomAlts: Ignorance

(30) Alice has more/less than 3 / at most/least 3 diamonds.

a.
$$O_{SubDomAlts}^{PS}$$
 $(|P \cap Q|) \in D$)
 $= |P \cap Q| \in D \land$
 $\neg |P \cap Q| \in A \land$
 $\neg |P \cap Q| \in B \land \dots$, for all $A, B, \dots \subset D$, $= \bot$
contradiction \checkmark
b. $O_{SubDomAlts}^{PS}$ $(\boxdot |P \cap Q| \in D)$
 $= \boxdot |P \cap Q| \in D \land$
 $\neg \boxdot |P \cap Q| \in B \land \dots$, for all $A, B, \dots \subset D$ ignorance \checkmark

- The only consistent O^{PS}_{SubDomAlts} parse yields ignorance.
 SMNs can only have an O^{PS}_{SubDomAlts} parse, so *(ignorance)
 CMNs can also have a parse without O^{PS}_{SubDomAlts}, so (ignorance).

Scalar implicatures vs. ignorance implicatures

(31) Alice has more than 2 /at least 3 diamonds. $O_{SubDomAlts}^{PS} \boxdot O_{ScalAlts} (|P \cap Q|) \in \{3, 4, \ldots\})$ $= \boxdot O_{ScalAlts} (|P \cap Q| \in \{3, 4, \dots\}) \land$ $\neg \boxdot (|P \cap Q| \in \{3\}) \land$ $\neg \boxdot (|P \cap Q| \in \{4, 7\}) \land \ldots$ $= \boxdot (|P \cap Q| \in \{3, 4, \ldots\} \land \neg |P \cap Q| \in \{4, \ldots\}) \land$ $\neg \boxdot (|P \cap Q| \in \{3\}) \land$ $\neg \boxdot (|P \cap Q| \in \{4, 7\}) \land \ldots$ $= \boxdot(|P \cap Q| \in \{3\}) \land \neg \boxdot (|P \cap Q| \in \{3\}) \land$ $\neg \boxdot (|P \cap O| \in \{4, 7\}) \land \ldots = \bot$ contradiction X PS

- * Prune offending SubDomAlts? That would violate $O_{SubDomAlts}^{[10]}$.
- ⋆ Prune offending ScalAlt? ✓

Implicatures from SubDomAlts: Negation

(32) Alice doesn't have more/less than three / *at most/least three diamonds.

a.
$$\neg O_{SubDomAlts}^{PS} (|P \cap Q| \in D)$$

b. $O_{SubDomAlts}^{PS} (\neg |P \cap Q| \in D)$
c. $O_{SubDomAlts}^{PS} (\boxdot \neg |P \cap Q| \in D)$



- * All $O_{SubDomAlts}^{PS}$ parses fail.
- SMNs cannot have a non-O^{PS}_{SubDomAlts} parse, so bad.
 CMNs can be parsed without O^{PS}_{SubDomAlts}, so okay.

Implicatures from SubDomAlts: AntCond/RestUniv



(33) $O_{SubDomAlts}^{PS}$ (Everyone who has at least 3 diamonds wins.)

Prejacent: $\forall x [\# \text{ di } x \text{ has } \in D \to \dots] \land \exists x [\# \text{ of di } x \text{ has } \in D]$ $\downarrow \qquad \qquad \uparrow$ SubDomAlt: $\forall x [\# \text{ di } x \text{ has } \in D' \to \dots] \land \exists x [\# \text{ of di } x \text{ has } \in D']$

- ★ SubDomAlts not entailed, so they must be false.
- ★ However, negating them leads to contradiction.
- ★ We can rescue the parse with ⊡:

 $(34) \boxdot \exists x [\# \text{ of di } x \text{ has } \in D] \land \neg \boxdot \exists x [\# \text{ of di } x \text{ has } \in D']$ $PS \text{ satisfied } \checkmark$

* Thus there *is* a consistent $O_{SubDomAlts}^{PS}$ parse for SMNs, which is why they are felicitous in this type of DE environments.

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The existential implicature of at most

[Alrenga, 2016]

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(35) LeBron scored at most 20 points (and it's even possible that
he didn't score any points at all).

O_{ScalAlts} (|P \cap Q| \in [[much]] (20))

= |P \cap Q| \in [[much]] (20) \land

\neg |P \cap Q| \in [[much]] (18) \land

\neg |P \cap Q| \in [[much]] (17) \land

...

\neg |P \cap Q| \in [[much]] (0)

existential implicature \checkmark
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★ Lower-bounding inference is a scalar implicature, which is why it is defeasible.

★ The same can be observed for *less than*.

 $\star\,$ Both follow if we assume CMNs and SMNs have scalar alternatives.

The 'not possible more' reading of *at most* under \Diamond

(36) a. You are allowed to drink at most one beer. b. $O_{ScalAlts} (O_{SubDomAlts}^{PS} (\diamond | P \cap Q | \in \llbracket \text{much} \rrbracket (1)))$ c. Prejacent: $O_{ExhSubDomAlts}^{PS} (\diamond | P \cap Q | \in \llbracket \text{much} \rrbracket (1))$ d. ScalAlts: $O_{ExhSubDomAlts}^{PS} (\diamond | P \cap Q | \in \llbracket \text{much} \rrbracket (m))$ e. Outcome: $\begin{array}{l} O_{ScalAlts} \left(O_{ExhSubDomAlts}^{PS} \left(\Diamond | P \cap Q | \in \llbracket \text{much} \rrbracket \left(1 \right) \right) \right) \\ = O_{ExhSubDomAlts}^{PS} \left(\Diamond | P \cap Q | \in \llbracket \text{much} \rrbracket \left(1 \right) \right) \land \\ \neg O_{SubDomAlts}^{PS} \left(\Diamond | P \cap Q | \in \llbracket \text{much} \rrbracket \left(2 \right) \right) \end{array}$ $= \Diamond |P \cap Q| \in \{0\} \land \Diamond |P \cap Q| \in \{1\} \land$ $\neg(\Diamond | P \cap Q| \in \{0\} \land \Diamond | P \cap Q| \in \{1\} \land \Diamond | P \cap Q| \in \{2\})$ not possible more \checkmark

★ This follows from a system where $O_{SubDomAlts}$ can apply to pre-exhaustified alternatives, where $O_{SubDomAlts}$ and $O_{ScalAlts}$ can be manipulated separately, and where $O_{SubDomAlts}$ can be part of the prejacent and the alternatives operated on by $O_{ScalAlts}$.

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Conclusion

★ A unified account of bare, comparative-modified, and superlative-modified numerals that

- captures more patterns than previous accounts; and

- derives them from

- truth conditions and alternatives obtained in a uniform way from the morphological pieces of BNs, CMNs, and SMNs, and

- general implicature calculation mechanisms, using general recipes for deriving scalar implicatures, ignorance effects, polarity sensitivity, or free choice behavior.

Open issues

★ How does superlativity in SMNs (at-sup) connect to superlativity in adjectives (sup)?

★ Why Proper Strengthening?

(At present it is a stipulation. We could replace it with a ban on vacuous exhaustification but I think in the general case that might be too strong. It however seems to be a necessary general assumption for items with a positive polarity behavior such as SMNs - [Spector, 2014, Nicolae, 2017]. Parametric choice? Is there any evidence of SMNs that are not PPIs?)

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Truth conditions, BNs

(37) 3 people quit.



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Truth conditions, CMNs

(38) More/less than 3 people quit.



Truth conditions, SMNs

(39) At most/least 3 people quit.



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